

Characterization of minimal cycle obstruction sets for partitionable planar graphs

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October 19, 2016

Background

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- Partition graphs that belong to some more general class?
- Partition the vertex-set into “stable-like” sets?
 - Each set induces a graph with bounded maximum degree.
 - Each set induces a graph with bounded component size.

A more general class

Hadwiger's Conjecture: The vertex-set of every K_{t+1} -minor-free graph can be partitioned into t stable sets.

Theorem: Let G be a K_{t+1} -minor free graph.

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- ② **(Edwards, Kang, Kim, Oum, Seymour)** For every k , there exists a K_{t+1} -minor-free graph that does not admit a partition into $t - 1$ parts where each induces a graph of maximum degree less than k .

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- 4 **(L., Oum)** $V(G)$ can be partitioned into 3 parts where each part induces a graph of bounded component size, if G has bounded maximum degree.

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- 5 For every k , there exists a planar graph with maximum degree 6 that does not admit a partition into 2 parts where each induces a graph of diameter less than k .

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- 2 For every k , there exists a planar graph such that for every partition into 2 parts, some part induces a graph of maximum degree at least k .

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- 3 **(Esperet, Joret)** For every k , there exists a planar graph that does not admit a partition into 3 parts each induces a graph of diameter less than k .

Planar graphs with some cycles forbidden

Let S be a set. A graph is S -free if it does not contain any subgraph isomorphic to some member of S .

Grötzsch's Theorem: Every $\{C_3\}$ -free planar graph is 3-colorable.

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Steinberg's conjecture: Every $\{C_4, C_5\}$ -free planar graph is 3-colorable.
(Disproved by **Cohen-Addad, Hebdige, Král', Li, Salgado.**)

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How about partitioning $V(G)$ into stable-like sets for planar graphs G with some cycles forbidden?

Improper coloring

A graph is (k_1, k_2, \dots, k_t) -colorable if its vertex-set can be partitioned into t sets X_1, X_2, \dots, X_t such that $G[X_i]$ has maximum degree at most k_i for every i .

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Theorem: Let G be a planar graph.

- 1 (4CT) G is $(0, 0, 0, 0)$ -colorable.
- 2 (Cowen, Cowen, Woodall) G is $(2, 2, 2)$ -colorable.

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Theorem: Let G be a planar graph.

- 1 **(4CT)** G is $(0, 0, 0, 0)$ -colorable.
- 2 **(Cowen, Cowen, Woodall)** G is $(2, 2, 2)$ -colorable.
- 3 For every k , some planar graph is not $(1, k, k)$ -colorable.

Theorem: Let G be a planar graph.

- 1 **(Grötzsch)** If G is $\{C_3\}$ -free, then G is $(0, 0, 0)$ -colorable.
- 2 If G is $\{C_3, C_4\}$ -free, then G is
 - $(1, 10)$ -colorable **(Choi, Choi, Jeong, Suh)**
 - $(2, 6)$ -colorable **(Borodin, Kostochka)**
 - $(3, 5)$ -colorable **(Choi, Raspaud)**
 - $(4, 4)$ -colorable **(Škrekovski)**

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A graph is $(0, *)$ -colorable if there exists M such that it is $(0, M)$ -colorable.

Theorem (Choi, L., Oum)

- 1 The minimal $(*, *)$ -obstruction sets are $\{C_4\}$ and the set of all odd cycles.
- 2 The minimal $(0, *)$ -obstruction sets are $\{C_3, C_4, C_6\}$ and the set of all odd cycles.
- 3 The minimal $(0, 0, *)$ -obstruction sets are $\{C_3\}$ and $\{C_4\}$.

Lemma

For every positive integer k and ℓ , there exist non- $(0, k)$ -colorable planar graphs G_1, G_2, G_3, G_4 such that

- *every cycle in G_1 has length 4 or $2\ell + 1$,*
- *every cycle in G_2 has length 3,*
- *every cycle in G_3 has length 6 or $4\ell + 1$, and*
- *every cycle in G_4 has length 6 or $4\ell + 3$.*

$(0, *)$ -obstruction

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- *every cycle in G_4 has length 6 or $4\ell + 3$.*

Lemma

*If S is a $(0, *)$ -obstruction set, then either S contains all odd cycles, or S contains $\{C_3, C_4, C_6\}$.*

Lemma

Let G be an S -free planar graph.

- 1 If S contains all odd cycle, then G is $(0, 0)$ -colorable.
- 2 If $S = \{C_3, C_4, C_6\}$, then G is $(0, 45)$ -colorable.

$(0, *)$ -obstruction

Lemma

Let G be an S -free planar graph.

- 1 If S contains all odd cycle, then G is $(0, 0)$ -colorable.
- 2 If $S = \{C_3, C_4, C_6\}$, then G is $(0, 45)$ -colorable.

Theorem

The minimal $(0, *)$ -obstruction sets are $\{C_3, C_4, C_6\}$ and the set of all odd cycles.

- Minimal obstruction sets for partitioning planar graphs into (2 or 3) graphs with bounded component size?
- Minimal obstruction sets for partitioning more general graphs into graphs with bounded maximum degree?
- Is it possible to partition every K_{t+1} -minor-free graph with no triangle into less than t graphs with bounded maximum degree/component size?
- Minimal obstruction sets for partitioning graphs into sparse graphs?

THANK YOU!